Proof of statement (H2): From (50)

$$J_{l+1,T}(k, z+1, t) = J_{l+1,T}(k, z, t+1). \qquad (55)$$

Thus, from (53), $J_{l+1,T}(k, z+1, t) \leq J_{l+1,T}(k, z, t)$. Then, (H2) follows from (18) and (54).

Proof of statement (H3): From (48) and (54)

$$J_{l+1,B}(k, z+1, t-1) \leq J_{l+1,B}(k, z, t).$$

From (55)

$$J_{l+1,T}(k, z+1, t-1) \leq J_{l+1,T}(k, z, t).$$

Then, (H3) follows from (18).

Thus, (H1), (H2), and (H3) hold for all l.

After taking limits as l goes to ∞ in (H1), it follows that the optimal policy is threshold type.

Now, we show that the algorithm in Fig. 10 obtains a threshold that minimizes the expected termination time for every $k \leq K$ and $z < Z$. Let $G^{\pi}(k, z)$ denote the expected time to terminate under a policy π after the kth transmission and the subsequent backoff, if z receivers are satisfied after k transmissions.

We show that for every $k \leq K - 1$ and z

$$G^*(k, z) = G^{\pi^*}(k, z). \qquad (56)$$

Since $G^*(k, z) = G^{\pi_2(K, Z)}(k, z)$, (56) proves the optimality of $\pi_2(K, Z)$.

Note that if $z \geq Z$, then $G^*(k, z) = G^{\pi^*}(k, z) = 0$ for every k. Thus, (56) follows. Henceforth, we consider $z < Z$.

Let $k = K - 1$. Clearly, π^* transmits when at least $Z - z$ unsatisfied receivers are ready. Thus, $G^{\pi^*}(K-1, z) = \frac{\overline{X}}{p_{Z-z}} + \overline{V}$ where $\{\widehat{p_u}\}$ are as defined in Fig. 10. Thus, (56) follows.

Now, we assume (56) for every $\widehat{k} > k$ and show (56) for k. Clearly

$$G^{\pi^*}(k, z) = \frac{\overline{X}}{\widehat{p}_{m_{k,z}}} + \overline{V} + \sum_{v=m_{k,z}}^{G-z-v} q_{m_{k,z},v}(z+v) G^{\pi^*}(k+1, z+v) \qquad (57)$$

where $\{q_{u,v}(z)\}$ are as defined in Fig. 10. Now, from Lemmas 3 and 5, $0 < m_{k,z} \leq Z - z$. Thus, from (57)

$$\begin{aligned} G^{\pi^*}(k, z) \\ \geq \min_{1 \leq u \leq Z-z} &\left\{ \frac{\overline{X}}{\widehat{p}_u} + \overline{V} + \sum_{v=u}^{G-z-v} q_{u,v}(z+v) G^{\pi^*}(k+1, z+v) \right\} \\ = \min_{1 \leq u \leq Z-z} &\left\{ \frac{\overline{X}}{\widehat{p}_u} + \overline{V} + \sum_{v=u}^{G-z-v} q_{u,v}(z+v) G^*(k+1, z+v) \right\} \\ &\text{(from induction hypothesis)}. \qquad (58) \end{aligned}$$

Clearly

$$G^{\pi^*}(k, z) \leq G^{\pi_2(K, Z)}(k, z) = G^*(k, z).$$

Thus, $G^{\pi^*}(k, z) = G^*(k, z)$. The result follows. □

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Unitary Space–Time Constellation Analysis: An Upper Bound for the Diversity

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Abstract—The diversity product and the diversity sum are two very important parameters for a good-performing unitary space-time constellation. A basic question is what the maximal diversity product (or sum) is. In this correspondence, we are going to derive general upper bounds on the diversity sum and the diversity product for unitary constellations of any dimension n and any size m using packing techniques on the compact Lie group $U(n)$.

Index Terms—Diversity product, diversity sum, multiple antennas, space–time coding, space–time constellations.

I. Introduction

Let A be a matrix with complex entries and A^* denote the conjugate transpose of A. Let $\| \cdot \|$ denote the Frobenius norm of a matrix, i.e.,

$$\|A\| = \sqrt{\mathrm{tr}(AA^*)}.$$

A square matrix A is called unitary if $A^*A = AA^* = I$, where I denotes the identity matrix. We denote by $U(n)$ the set of all $n \times n$

unitary matrices. $U(n)$ is a real algebraic variety and a smooth manifold of real dimension n^2. For the purpose of this correspondence, a unitary space–time constellation (or code) $\mathcal{V}$ is simply a finite subset of $U(n)$

$$\mathcal{V} = \{A_1, A_2, \ldots, A_m\} \subset U(n).$$

We say $\mathcal{V}$ has dimension n and size m. Unitary space–time codes have been intensely studied in recent years and we refer the interested readers to [1], [12], [13], [16], and the references of these papers. The readers will find the motivation and engineering applications of such kind of codes. The quality of a unitary space–time code is governed by two important parameters, the diversity product and the diversity sum.

Definition 1.1: The *diversity product* [12] of a unitary space–time code $\mathcal{V}$ is defined through

$$\prod \mathcal{V} := \frac{1}{2} \min \left\{ |\det(A - B)|^{\frac{1}{n}} \,|\, A, B \in \mathcal{V}, A \neq B \right\}.$$

The *diversity sum* [14] is defined as

$$\sum \mathcal{V} := \frac{1}{2\sqrt{n}} \min\{\|A - B\| \,|\, A, B \in \mathcal{V}, A \neq B\}.$$

$\mathcal{V}$ is called fully diverse if $\prod \mathcal{V} > 0$. As explained in [8], a space–time code with large diversity sum tends to perform well at low signal-to-noise ratios whereas a code with a large diversity product tends to perform well at high signal-to-noise ratios. A major coding design problem is the construction of unitary space–time codes where the diversity sum (or product) is optimal or near optimal inside the set of all the space–time codes with the same parameters n, m. We would like to remark that for every positive integer n and m, a Haar distributed random space–time code is fully diverse with probability 1.

The purpose of this correspondence is to derive for n and m tight upper bounds for the diversity product $\prod \mathcal{V}$ and the diversity sum $\sum \mathcal{V}$. When $n = 1$, trivially $|\det(A - B)| = \|A - B\|$; and it follows that $\sum \mathcal{V} = \prod \mathcal{V}$ in this situation. For any unitary space–time code $\mathcal{V}$, we have $\prod \mathcal{V} \leq \sum \mathcal{V}$ (see [14]). Thus, by having an upper bound for $\sum \mathcal{V}$ we immediately also have an upper bound for $\prod \mathcal{V}$.

Of course it would be desirable to know for every n and m what the largest possible value of $\sum \mathcal{V}$ is. This is the motivation of the following definition.

Definition 1.2: Let $\Delta(n, m)$ be the infimum of all numbers such that for every unitary space–time code $\mathcal{V}$ of dimension n and size m, one has

$$\sum \mathcal{V} \leq \Delta(n, m).$$

Remark 1.3: As pointed out by Liang and Xia [14], there exists a constellation $\mathcal{V}$ of dimension n and size m with $\sum \mathcal{V} = \Delta(n, m)$. This is due to the fact that $U(n)^m$ is a compact manifold.

The exact values of $\Delta(n, m)$ are only known in very few special cases. In the case $n = 1$, one checks that $\Delta(1, m) = \sin \frac{\pi}{m}$ for $m \geq 2$. When $n \geq 2$ and $m = 3$, one has $\Delta(n, 3) = \frac{\sqrt{3}}{2}$. When $m = 2$, we have $\Delta(n, 2) = 1$ for $n \geq 2$. For $n = 2$, the values shown in the table at the bottom of the page were computed in [14].

Liang and Xia [14] observed the connection between a unitary constellation and an Euclidean sphere code and derived an upper bound for two-dimensional unitary constellations which is very tight when $m \leq 100$. In this correspondence, we present a new general upper bound for $\Delta(n, m)$ for every dimension n and every size m while improving certain results in [14]. To the best of our knowledge, the new upper bounds we derived are tighter than any previously published bounds as soon as m is sufficiently large. Independently from this paper, Henkel [11] derived recently also upper bounds for the problem we study.

II. UPPER BOUND ANALYSIS

In this section, we are going to study the packing problem on $U(n)$ and derive three upper bounds for the numbers $\Delta(n, m)$. All the resulting bounds are derived by differential geometric means and all bounds can be viewed as certain sphere-packing bounds.

From a differential geometry point of view, we can view $U(n)$ as an n^2-dimensional compact Lie group. $U(n)$ is also naturally a submanifold of the Euclidean space $\mathbb{R}^{2n^2}$. In this way, $U(n)$ will have the induced geometry of the standard Euclidean geometry of $\mathbb{R}^{2n^2}$. Finally, there is a third way to see $U(n)$ as a submanifold of another Riemannian manifold $S(n)$ and we will say more later.

In the sequel, we will employ standard techniques from Riemannian geometry in order to derive upper bounds for the maximal diversity of a constellation. Some standard textbooks on the subjects of Riemannian geometry, differentiable manifolds and Lie groups are, e.g., [3], [10], and [15].

The basic strategy for computing the upper bounds for $\Delta(n, m)$ is as follows. Given a unitary space–time code $\mathcal{V} = \{A_1, A_2, \ldots, A_m\}$, around each matrix A_j we can choose a neighborhood $N_r(A_j)$ with radius r (the radius will be specified later). Let $V_j = V(N_r(A_j))$ be the volume of the neighborhood $N_r(A_j)$. If all the neighborhoods are nonoverlapping, then necessarily we will have

$$\sum_{j=1}^{m} V_j \leq V(U(n))$$

where $V(U(n))$ denotes the total volume of unitary group $U(n)$. This inequality in turn will result in an upper bound for the numbers $\Delta(n, m)$. By employing different metrics (Euclidean or Riemannian) and by considering different embeddings of $U(n)$, we derive three different upper bounds for $\Delta(n, m)$.

Let $\mathcal{M}_1$ be the manifold consisting of all the $n \times n$ Hermitian matrices, i.e.,

$$\mathcal{M}_1 = \{H | H = H^*\}.$$

$\mathcal{M}_1$ has dimension n^2 and can be viewed isometrically (see, e.g., [3, p. 189]) as Euclidean space $\mathbb{R}^{n^2}$. Assume that $H = (H_{jk})$ and assume that $H_{jk} = x_{jk} + iy_{jk}$. We use (dH) to denote the volume element (see, e.g., [10, p. 216]) of $\mathcal{M}_1$, where

$$(dH) = \left(\frac{i}{2}\right)^{n(n-1)/2} \bigwedge_{l=1}^{n} dH_{ll} \bigwedge_{j<k} dH_{jk} \bigwedge_{j<k} d\bar{H}_{jk}$$

$$= \bigwedge_{l=1}^{n} dx_{ll} \bigwedge_{j<k} dx_{jk} \bigwedge_{j<k} dy_{jk}. \qquad (2.1)$$

m	2	3	4	5	6	7	8	9	10 through 16
$\Delta(2, m)$	1	$\frac{1}{2}\sqrt{3}$	$\frac{1}{3}\sqrt{6}$	$\frac{1}{4}\sqrt{10}$	$\frac{1}{5}\sqrt{15}$	$\frac{1}{6}\sqrt{21}$	$\frac{1}{7}\sqrt{28}$	$\frac{1}{8}\sqrt{36}$	$\frac{1}{2}\sqrt{2}$

With a small abuse of the notation, one can check that the volume element of $\mathcal{M}_2$, the manifold consisting of all the $n \times n$ skew-Hermitian matrices, can be written as

$$(dH) = \left(\frac{i}{2}\right)^{n(n-1)/2} \left(\frac{1}{i}\right)^n \bigwedge_{l=1}^n dH_{ll} \bigwedge_{j<k} dH_{jk} \bigwedge_{j<k} d\bar{H}_{jk}$$

$$= \bigwedge_{l=1}^n dy_{ll} \bigwedge_{j<k} dx_{jk} \bigwedge_{j<k} dy_{jk}. \tag{2.2}$$

For a unitary matrix U, if we differentiate $U^*U = I$, we will have

$$U^* dU + dU^* U = 0.$$

Therefore, $U^* dU$ is skew-Hermitian.

A differential form on a Lie group is left-invariant if it coincides with the differential form induced by any left-multiplication, and is right-invariant if it coincides with the differential form induced by any right-multiplication. We say a differential form on a Lie group is bi-invariant if it is left-invariant and right-invariant at the same time. For a compact connected Lie group, left-invariance (or right-invariance) implies bi-invariance [15]. The following lemma will characterize the volume element of $U(n)$.

Lemma 2.1: The volume element of $U(n)$ induced by the Euclidean space $\mathbb{R}^{2n^2}$ is bi-invariant and the volume element can be written as $(U^* dU)$ up to a scalar constant.

Proof: The bi-invariance comes from the orthonormality of $U(n)$. $(U^* dU)$ is left-invariant according to the definition. Indeed, for a fixed yet arbitrary unitary matrix V

$$(VU)^* d(VU) = U^* V^* V dU = U^* dU.$$

Since $U(n)$ is a compact Lie group, $(U^* dU)$ is also right-invariant. Because the bi-invariant n^2 differential forms are unique up to a scalar, one concludes that the volume element can be written as $(U^* dU)$. Because the bi-invariant n^2 differential forms are unique up to a scalar, one concludes that the volume element can be written as $(U^* dU)$. $\square$

The following theorem will represent the volume element of $U(n)$ in another way. One will see that it is closely related to the eigenvalues of unitary matrices.

Theorem 2.2: Consider the eigenvalue decomposition of a unitary matrix Θ

$$\Theta = U \operatorname{diag}\left(e^{i\theta_1}, e^{i\theta_2}, \ldots, e^{i\theta_n}\right) U^*. \tag{2.3}$$

In order to make this decomposition unique, we assume that $U \in \tilde{U}(n)$, where $\tilde{U}(n)$ denote the set of all the unitary matrices with non-negative real diagonal elements and that $\theta_1 > \theta_2 > \cdots > \theta_n$. Then we will have

$$(\Theta^* d\Theta) = \prod_{j<k} \left| e^{i\theta_j} - e^{i\theta_k} \right|^2 d\theta_1 \wedge d\theta_2$$

$$\wedge \cdots \wedge d\theta_n \wedge (U^* dU - \operatorname{diag}(U^* dU)). \tag{2.4}$$

Proof: Let $D = \operatorname{diag}\left(e^{i\theta_1}, e^{i\theta_2}, \ldots, e^{i\theta_n}\right)$ and take the differential of (2.3)

$$d\Theta = dU D U^* + U dD U^* + U D dU^*.$$

It follows that

$$\Theta^* d\Theta = U D^* U^* dU D U^* + U D^* dD U^* + U dU^*$$
$$= U(D^* U^* dU D + D^* dD) U^* + U dU^*.$$

Due to the right-invariance of the volume element in $U(n)$, it follows that

$$(\Theta^* d\Theta) = (U^* \Theta^* d\Theta U)$$
$$= (D^* U^* dU D - U^* dU + i \operatorname{diag}(d\theta_1, d\theta_2, \ldots, d\theta_n)).$$

Note that $(D^* U^* dU D - U^* dU)_{jk} = (e^{i\theta_j} - e^{i\theta_k}) U_{jk}$, therefore, the diagonal elements of $D^* U^* dU D - U^* dU$ are all zeros and the off-diagonal elements are scaled version of the ones of $U^* dU$. According to (2.2), the claim in the theorem follows. $\square$

The following theorem calculates the volume of a small neighborhood with Euclidean distance r. Because of the homogeneity of $U(n)$, the center of this small "ball" is chosen to be I without loss of generality. For a unitary matrix U, we assume $e^{i\theta_j}$'s are its eigenvalues, i.e., $U \sim \operatorname{diag}\left(e^{i\theta_1}, e^{i\theta_2}, \ldots, e^{i\theta_n}\right)$ (here $\sim$ means similar). For a fixed unitary matrix A, let

$$U_r^E(n, A) = \{U \in U(n) | \|U - A\| \leq r\}.$$

Again, because of the homogeneity of $U(n)$, $V(U_r^E(n, A))$ does not depend on the choice of A. In the sequel, $V(U_r^E(n))$ will be used to denote $V(U_r^E(n, A))$ for any unitary matrix A. Let $S(n)$ denote a $2n^2 - 1$-dimensional sphere centered at the origin with radius $\sqrt{n}$, i.e.,

$$S(n) = \left\{(x_1, x_2, \ldots, x_{2n^2}) | x_1^2 + x_2^2 + \cdots + x_{2n^2}^2 = n\right\}.$$

Apparently $U(n)$ is a submanifold of $S(n)$. For a particular point $S_0 \in S(n)$, let

$$S_r(n, S_0) = \{S \in S(n) | \|S - S_0\| \leq r\}.$$

Theorem 2.3: Let

$$D_1 = \{(\theta_1, \theta_2, \ldots, \theta_n) | -\pi \leq \theta_j < \pi, \text{ for } j = 1, 2, \ldots, n\} \tag{2.5}$$

and

$$D_2 = \left\{(\theta_1, \theta_2, \ldots, \theta_n) | \sum_{j=1}^n \sin^2 \frac{\theta_j}{2} \leq \frac{r^2}{4}\right\} \tag{2.6}$$

then

$$V(U_r^E(n)) = \frac{\iint_{D_1 \cap D_2} \prod_{j<k} \left| e^{i\theta_j} - e^{i\theta_k} \right|^2 d\theta_1 d\theta_2 \cdots d\theta_n}{\iint_{D_1} \prod_{j<k} \left| e^{i\theta_j} - e^{i\theta_k} \right|^2 d\theta_1 d\theta_2 \cdots d\theta_n} V(U(n)). \tag{2.7}$$

Proof: Note that $\|I - U\|_2 \leq r$ is equivalent to

$$\sum_{j=1}^n \sin^2 \frac{\theta_j}{2} \leq \frac{r^2}{4}.$$

For a given unitary matrix Θ, the eigenvalue decomposition $\Theta = U^* \operatorname{diag}\left(e^{i\theta_1}, e^{i\theta_2}, \ldots, e^{i\theta_n}\right) U$ is unique if θ_j's are strictly ordered. So if we take the integral of formula (2.4) over the integration region disregarding the order of θ_j's, we will obtain $n!$ times the volume of $V(U_r^E(n))$. Thus, the volume of $U_r^E(n)$ will be

$$V(U_r^E(n)) = \frac{1}{n!} \int \int_{D_1 \cap D_2} \prod_{j<k} \left| e^{i\theta_j} - e^{i\theta_k} \right|^2 d\theta_1 d\theta_2$$

$$\cdots d\theta_n \int \int_{\tilde{U}(n)} (U^* dU - \operatorname{diag}(U^* dU)).$$

Using the same argument, we will derive the volume of $U(n)$

$$V(U(n)) = \frac{1}{n!} \int\int_{D_1} \prod_{j<k} |e^{i\theta_j} - e^{i\theta_k}|^2 d\theta_1 d\theta_2$$

$$\cdots d\theta_n \int\int_{\tilde{U}(n)} (U^* dU - \mathrm{diag}\,(U^* dU)).$$

Comparing the two derived volume formulas, the claim in the theorem follows. $\square$

There are several approaches to derive upper bounds for the diversity sum. The first approach considers $U(n)$ as a submanifold of $S(n)$, then chooses the nonoverlapping neighborhoods to be small balls with radius r (with regard to the Euclidean distance). This will result in the first upper bound (B1) which we derive in this correspondence.

Theorem 2.4: Let D_1 and D_2 be defined as in (2.5) and (2.6). Assume $r_0^E = r_0^E(n, m)$ is the solution to the following equation (with variable r):

$$m \int\int_{D_1 \cap D_2} \prod_{j<k} |e^{i\theta_j} - e^{i\theta_k}|^2 d\theta_1 d\theta_2 \cdots d\theta_n$$

$$= \int\int_{D_1} \prod_{j<k} |e^{i\theta_j} - e^{i\theta_k}|^2 d\theta_1 d\theta_2 \cdots d\theta_n \quad (2.8)$$

then

$$\Delta(n, m) \le \sqrt{\frac{(r_0^E)^2}{n} - \frac{(r_0^E)^4}{4n^2}}. \quad (B1)$$

Proof: For a fixed yet arbitrary unitary constellation $\mathcal{V} = \{A_1, A_2, \ldots, A_m\}$, consider m small nonoverlapping neighborhoods $S_r(n, A_j)$ in $S(n)$. We can increase r such that there exist l, k such that $S_r(n, A_l)$ and $S_r(n, A_k)$ are tangent to each other. Apparently

$$U_r^E(n, A_j) = S_r(n, A_j) \cap U(n)$$

for any j. Since $S_r(n, A_j)$'s are nonoverlapping, we conclude that $U_r^E(n, A_j)$'s are nonoverlapping. Therefore, we have

$$\sum_{j=1}^m V\left(U_r^E(n, A_j)\right) \le V(U(n))$$

that is,

$$m V\left(U_r^E(n)\right) \le V(U(n)).$$

One can check that $V\left(U_r^E(n)\right)$ is an increasing function of r, so any r satisfying the above inequality will be less than the solution to the equality

$$m V(U_r^E(n)) = V(U(n))$$

which is essentially (2.8). So we conclude that $r \le r_0^E$.

Note that any two points $S_0, S_1 \in S(n)$ with two nonoverlapping neighborhoods $S_r(n, S_0)$ and $S_r(n, S_1)$ will have distance

$$\|S_0 - S_1\| \ge 2\sqrt{r^2 - r^4/(4n)}$$

where the equality holds only if $S_r(n, S_0)$ and $S_r(n, S_1)$ are tangent to each other. Apply the argument to A_j's and note that A_l and A_k are the closest pair of points with $\|A_l - A_k\| = 2\sqrt{r^2 - r^4/(4n)}$, we reach the conclusion of the theorem. $\square$

Remark 2.5: For the right-hand side of (2.8), by the Weyl denominator formula [6] one can replace

$$\int\int_{D_1} \prod_{j<k} |e^{i\theta_j} - e^{i\theta_k}|^2 d\theta_1 d\theta_2 \cdots d\theta_n$$

with $(2\pi)^n n!$. However, we are not able to obtain a closed-form expression for the left-hand side of (2.8), and it seems that we need to compute the following:

$$\int\int_{D_1 \cap D_2} \prod_{j<k} |e^{i\theta_j} - e^{i\theta_k}|^2 d\theta_1 d\theta_2 \cdots d\theta_n.$$

Hassibi suggested that the preceding expression can be simplified to facilitate the computation: first note that for any single variable function $f(\theta)$, we have

$$\int\int_{D_1} \prod_{j<k} |e^{i\theta_j} - e^{i\theta_k}|^2 \prod_{l=1}^n f(\theta_l) d\theta_1 d\theta_2 \cdots d\theta_n = n! \det(F)$$

$$(2.9)$$

where F is a Hankel matrix

$$F = \int_{-\pi}^{\pi} f(\theta) \begin{bmatrix} 1 \\ e^{i\theta} \\ \vdots \\ e^{i(n-1)\theta} \end{bmatrix} \begin{bmatrix} 1 & e^{-i\theta} & \cdots & e^{-i(n-1)\theta} \end{bmatrix} d\theta.$$

$$(2.10)$$

The proof of (2.9) is parallel to that of Lemma 1 in [9], and thus omitted. Next let $u(t)$ denote the rectangular function: $u(t)$ is equal to 0 if $t < 0$ or $t > 1$, to 1 if $0 \le t \le 1$. By the Fourier transform formula, we know that

$$u(t) = \frac{1}{2\pi i} \int_{-\infty}^{\infty} \frac{i}{w} (e^{-iw} - 1) e^{itw} dw.$$

Letting $\mathbf{1}_{D_2}$ denote the indicator function of D_2, we have

$$\mathbf{1}_{D_2} = \frac{1}{2\pi} \int_{-\infty}^{\infty} \frac{1}{w} (e^{-iw} - 1) e^{i\left(4\sin^2 \frac{\theta_1}{2} + \cdots + 4\sin^2 \frac{\theta_n}{2}\right)/r^2 w} dw.$$

Thus,

$$\int\int_{D_1 \cap D_2} \prod_{j<k} |e^{i\theta_j} - e^{i\theta_k}|^2 d\theta_1 d\theta_2 \cdots d\theta_n$$

$$= \int\int_{D_1} \prod_{j<k} |e^{i\theta_j} - e^{i\theta_k}|^2 \mathbf{1}_{D_2} d\theta_1 d\theta_2 \cdots d\theta_n$$

$$= \frac{1}{2\pi} \int \frac{1}{w} (e^{-iw} - 1)$$

$$\times \int\int_{D_1} \prod_{j<k} |e^{i\theta_j} - e^{i\theta_k}|^2 \prod_{l=1}^n f(\theta_l) d\theta_1 d\theta_2 \cdots d\theta_n dw;$$

here $f(\theta) = e^{i\frac{4\sin^2 \theta}{r^2} w}$. Applying (2.9), we obtain

$$\int\int_{D_1 \cap D_2} \prod_{j<k} |e^{i\theta_j} - e^{i\theta_k}|^2 d\theta_1 d\theta_2 \cdots d\theta_n$$

$$= \frac{n!}{2\pi} \int \frac{1}{w} (e^{-iw} - 1) \det(F) dw,$$

where F is defined in (2.10). In other words, the multiple-fold integral of the left-hand side of (2.8) can be transformed to a single-fold integral, which will facilitate the numerical computation for large n.

For a fixed $S_0 \in S(n)$, consider $S_r(n, S_0) \subset S(n)$. Let $\tau = \tau(n, r)$ denote the maximal number τ such that $S_r(n, S_1)$, $S_r(n, S_2), \ldots, S_r(n, S_\tau)$ are nonoverlapping and $S_r(n, S_j)$ is tangent to $S_r(n, S_0)$ for $j = 1, 2, \ldots, \tau$. One checks that $\tau(n, r)$

does not depend on the choice of S_0. In this sense, $\tau(n,r)$ can be viewed as generalized kissing number [4] on an Euclidean sphere. For a fixed n-dimensional unitary constellation $\mathcal{V} = \{A_1, A_2, \ldots, A_m\}$, let $r(\mathcal{V})$ denote the maximal radius r such that $S_r(n, A_1), S_r(n, A_2), \ldots, S_r(n, A_m)$ are nonoverlapping. Let $r_{\mathrm{opt}} = r_{\mathrm{opt}}(n, m)$ denote the maximal $r(\mathcal{V})$ over all possible n-dimensional unitary constellation $\mathcal{V}$ with cardinality m. One checks $\Delta(n, m) = r_{\mathrm{opt}}(n, m)/\sqrt{2n}$. The following theorem and corollary give a lower bound for the optimal diversity sum $\Delta(n, m)$.

Theorem 2.6: Let D_1 be defined as in (2.5) and assume that $r_0^E = r_0^E(n, m)$ is the solution to (2.8). Let

$$\tilde{D}_2 = \left\{ (\theta_1, \theta_2, \ldots, \theta_n) \mid \sum_{j=1}^n \sin^2 \frac{\theta_j}{2} \leq \frac{(r_0^E)^2}{4} \right\}$$

and let

$$D_3 = \left\{ (\theta_1, \theta_2, \ldots, \theta_n) \mid \sum_{j=1}^n \sin^2 \frac{\theta_j}{2} \leq r_{\mathrm{opt}}(n, m)^2 \right.$$
$$\left. - r_{\mathrm{opt}}(n, m)^4 / (4n) \right\}.$$

Then

$$\iint_{D_1 \cap \tilde{D}_2} \prod_{j<k} |e^{i\theta_j} - e^{i\theta_k}|^2 d\theta_1 d\theta_2 \cdots d\theta_n$$
$$\leq (\tau(n, r_{\mathrm{opt}}(n, m)) + 1)$$
$$\times \iint_{D_1 \cap D_3} \prod_{j<k} |e^{i\theta_j} - e^{i\theta_k}|^2 d\theta_1 d\theta_2 \cdots d\theta_n.$$

Proof: According to the derivation of r_0^E, we have

$$m \iint_{D_1 \cap \tilde{D}_2} \prod_{j<k} |e^{i\theta_j} - e^{i\theta_k}|^2 d\theta_1 d\theta_2 \cdots d\theta_n$$
$$= \iint_{D_1} \prod_{j<k} |e^{i\theta_j} - e^{i\theta_k}|^2 d\theta_1 d\theta_2 \cdots d\theta_n. \quad (2.11)$$

Assume that $\mathcal{V} = \{A_1, A_2, \ldots, A_m\}$ is an n-dimensional unitary constellation reaching $r_{\mathrm{opt}}(n, m)$, i.e., $r(\mathcal{V}) = r_{\mathrm{opt}}(n, m)$. For simplicity, let $r = r(\mathcal{V})$. Let m' denote the maximal number such that $S_r(n, A_1), S_r(n, A_2), \ldots, S_r(n, A_m), \ldots, S_r(n, A_{m'})$ are nonoverlapping. Let $r_1 = 2\sqrt{r^2 - r^2/(4n)}$, we claim that

$$U(n) \subset \bigcup_{j=1}^{m'} U_{r_1}^E(n, A_j).$$

Otherwise, suppose there is a unitary matrix $A_0 \notin \bigcup_{j=1}^{m'} U_{r_1}^E(n, A_j)$, then $\|A_0 - A_j\| > r_1$ (see Theorem 2.4). Thus, $S_r(n, A_0)$ does not intersect with $S_r(n, A_j)$ for $j = 1, 2, \ldots, m'$. Therefore, one can find $m' + 1$ small balls with radius r which are nonoverlapping. This contradicts the maximality of m'. Thus, we have

$$\sum_{j=1}^{m'} V(U_{r_1}^E(n, A_j)) \geq V(U(n))$$

that is,

$$m' \iint_{D_1 \cap D_3} \prod_{j<k} |e^{i\theta_j} - e^{i\theta_k}|^2 d\theta_1 d\theta_2 \cdots d\theta_n$$
$$\geq \iint_{D_1} \prod_{j<k} |e^{i\theta_j} - e^{i\theta_k}|^2 d\theta_1 d\theta_2 \cdots d\theta_n. \quad (2.12)$$

We further claim that

$$m' \leq (m-1)(\tau(n, r) + 1). \quad (2.13)$$

By contradiction assume that $m' \geq (m-1)(\tau(n, r) + 1) + 1$. Let

$$\mathrm{tang}(j) = \{l \mid 1 \leq l \leq m', S_r(n, A_l) \text{ tangent to } S_r(n, A_j)\}.$$

According to the definition of $\tau(n, r)$, we know the cardinality of $\mathrm{tang}(j)$ is less than $\tau(n, r)$. We first pick j_1 from $\{0, 1, \ldots, m'\}$, then pick j_2 from $\{0, 1, \ldots, m'\} - \mathrm{tang}(j_1)$. And we continue this process by always picking j_{k+1} from

$$\{0, 1, \ldots, m'\} - \bigcup_{l=1}^k \mathrm{tang}(j_l).$$

Since the cardinality of the above set is strictly greater than 0 when $k \leq m - 1$, we can pick $j_1, j_2, \ldots, j_m$ from the index set $\{1, 2, \ldots, m'\}$ such that $S_r(n, A_{j_1}), S_r(n, A_{j_2}), \ldots, S_r(n, A_{j_m})$ are nonoverlapping and every two of them are not tangent to each other. Then we can find a small enough real number $\varepsilon > 0$ and increase the radius r to $r + \varepsilon$ such that

$$S_{r+\varepsilon}(n, A_{j_1}), S_{r+\varepsilon}(n, A_{j_2}), \ldots, S_{r+\varepsilon}(n, A_{j_m})$$

are still nonoverlapping. However, this contradicts the maximality of $r = r_{\mathrm{opt}}(n, m)$. The combination of the three formulas (2.11)–(2.13) will lead to

$$\frac{\iint_{D_1} \prod_{j<k} |e^{i\theta_j} - e^{i\theta_k}|^2 d\theta_1 d\theta_2 \cdots d\theta_n}{\iint_{D_1 \cap D_3} \prod_{j<k} |e^{i\theta_j} - e^{i\theta_k}|^2 d\theta_1 d\theta_2 \cdots d\theta_n}$$
$$\leq \left(\frac{\iint_{D_1} \prod_{j<k} |e^{i\theta_j} - e^{i\theta_k}|^2 d\theta_1 d\theta_2 \cdots d\theta_n}{\iint_{D_1 \cap \tilde{D}_2} \prod_{j<k} |e^{i\theta_j} - e^{i\theta_k}|^2 d\theta_1 d\theta_2 \cdots d\theta_n} - 1 \right)$$
$$\times (\tau(n, r) + 1). \quad (2.14)$$

Note that the preceding inequality is, in fact, stronger than the claim in the theorem. We can reach the conclusion of the theorem by relaxing the right-hand side of the inequality (by ignoring -1). $\square$

Corollary 2.7: When $m \to \infty$, asymptotically we have

$$\Delta(n, m) \geq 2\sqrt{n} r_0^E(n, m) \frac{1}{2} (\tau(2n^2 - 1) + 1)^{-1/n^2}.$$

Proof: We only sketch the idea of the proof. Intuitively $U_r^E(n, A_0)$ looks more "flat" when $m \to \infty$ (consequently, $r \to 0$), so $V(U_r^E(n, A_0))$ can be approximated by the volume of $U_r^E(n, A_0)$'s projection to the tangent space of $U(n)$ at A_0

$$\iint_{D_1 \cap \tilde{D}_2} \prod_{j<k} |e^{i\theta_j} - e^{i\theta_k}|^2 d\theta_1 d\theta_2 \cdots d\theta_n \sim C(r_0^E)^{n^2}$$

for some constant C. The same argument will lead to

$$\iint_{D_1 \cap D_3} \prod_{j<k} |e^{i\theta_j} - e^{i\theta_k}|^2 d\theta_1 d\theta_2 \cdots d\theta_n \sim C(2r_{\mathrm{opt}})^{n^2}$$

for the same constant C. For any fixed n, $\tau(n, r)$ will approach the standard kissing number in Euclidean space $\tau(2n^2 - 1)$ when r goes to zero. Combining the three approximations, we reach the claim according to the previous theorem. $\square$

$U(n)$ is a compact Lie group equipped with a Riemannian metric. Given two points $A_0, A_1 \in U(n)$, one can always find a geodesic $\gamma(t)$ (mapping from $[0, 1]$ to $U(n)$) which will connect these two points, i.e., $\gamma(0) = A_0$ and $\gamma(1) = A_1$. Recall that the Euclidean distance of A_0

and A_1 is defined to be $\|A_0 - A_1\|$. We further define the Riemannian distance between A_0 and A_1 to be

$$\text{dist } (A_0, A_1) = \int_0^1 \|\gamma'(t)\| dt.$$

As a Lie group $U(n)$ is homogeneous. In particular one has that

$$\text{dist } (A_0, A_1) = \text{dist } (U A_0, U A_1) = \text{dist } (A_0 U, A_1 U)$$

for any $U \in U(n)$. The following theorem utilizes the homogeneity and the relationship between the Riemannian distance and Euclidean distance to derive another upper bound for the diversity sum in general and it is the base of the second approach.

Theorem 2.8: Let $f(\cdot)$ and $g(\cdot)$ be two fixed monotone increasing real functions. If

$$g(\|A_0 - A_1\|) \le \text{dist } (A_0, A_1) \le f(\|A_0 - A_1\|)$$

for any two unitary matrices A_0 and A_1, then

$$\Delta(n, m) \le g^{-1} \left(2f \left(r_0^E(n, m) \right) \right) / (2\sqrt{n}).$$

Proof: For a fixed unitary constellation $\mathcal{V} = \{A_1, A_2, \ldots, A_m\}$, consider

$$U_r^E(n, A_1), U_r^E(n, A_2), \ldots, U_r^E(n, A_m)$$

for $r > 0$. We can increase r until there exist j and k such that $U_r^E(n, A_j)$ and $U_r^E(n, A_k)$ are tangent to each other at a point A_0. As examined in Theorem 2.4, one can conclude that $r \le r_0^E(n, m)$. Accordingly, we have

$$\begin{aligned} \text{dist } (A_j, A_k) &\le \text{dist } (A_j, A_0) + \text{dist } (A_k, A_0) \\ &\le f(\|A_j - A_0\|) + f(\|A_k - A_0\|) \\ &= 2f(r) \le 2f \left(r_0^E(n, m) \right). \end{aligned}$$

On the other hand, since g is monotonically increasing one has

$$\|A_j - A_k\| \le g^{-1}(\text{dist } (A_j, A_k)).$$

The combination of the above two inequalities will lead to

$$\|A_j - A_k\| \le g^{-1} \left(2f(r_0^E(n, m)) \right).$$

Immediately, we will have

$$\sum \mathcal{V} \le g^{-1} \left(2f \left(r_0^E(n, m) \right) \right) / (2\sqrt{n}).$$

Since $\mathcal{V}$ is an arbitrary unitary constellation, the claim in the theorem follows. $\square$

Based on the above theorem, the following corollary gives the second upper bound (B2).

Corollary 2.9: For a real number r, let $\lfloor r \rfloor$ denote the greatest integer less than or equal to r, then we get (B2) at the bottom of the page.

Proof: Consider I and another point

$$U = V \text{diag } (e^{i\theta_1}, e^{i\theta_2}, \ldots, e^{i\theta_n}) V^*$$

where $-\pi \le \theta_j < \pi$. It is known that [5] the geodesic from I to U can be parameterized by

$$\gamma(t) = V \text{diag } (e^{i\theta_1 t}, e^{i\theta_2 t}, \ldots, e^{i\theta_n t}) V^*$$

where $0 \le t \le 1$. The Riemannian distance from I to U is

$$\text{dist } (I, U) = \sqrt{\theta_1^2 + \theta_2^2 + \cdots + \theta_n^2}.$$

We want to derive $g(\cdot), f(\cdot)$ as in Theorem 2.8. Suppose the Euclidean distance between I and U is r, i.e.,

$$\sin^2 \frac{\theta_1}{2} + \sin^2 \frac{\theta_2}{2} + \cdots + \sin^2 \frac{\theta_n}{2} = r^2/4.$$

After substituting with $x_j = \sin^2 \theta_j / 2$ and denoting $G(x) = \arcsin^2 \sqrt{x}$, we convert the above problem to the following optimization problem.

Find the minimum and maximum of the function

$$\begin{aligned} F(x_1, x_2, \ldots, x_n) &= \theta_1^2 + \theta_2^2 + \cdots + \theta_n^2 \\ &= 4(G(x_1) + G(x_2) + \cdots + G(x_n)) \end{aligned}$$

with the constraints $x_1 + x_2 + \cdots + x_n = r^2/4$ and $0 \le x_j \le 1$ for $j = 1, 2, \ldots, n$. Since $G(x)$ is a convex function on $[0, 1]$, we derive the lower bound of $F(x_1, x_2, \ldots, x_n)$

$$4n \arcsin^2(r/(2\sqrt{n})) \le F(x_1, x_2, \ldots, x_n). \qquad (2.15)$$

In the sequel, we are going to calculate the upper bound of $F(x_1, x_2, \ldots, x_n)$. Without loss of generality, we assume $0 \le x_1 \le x_2 \le \cdots \le x_n \le 1$. Let $k = \lfloor r^2/4 \rfloor$ and $\alpha = r^2/4 - k$, we claim that $F(x_1, x_2, \ldots, x_n)$ will reach its maximum when

$$x_j = \begin{cases} 1, & \le j \le n - k - 1 \\ \alpha, & j = n - k \\ 1, & n - k + 1 \le j \le n. \end{cases}$$

Suppose by contradiction that F reaches its maximum at $(x_1, x_2, \ldots, x_n)$ with $x_1 > 0$. Now from

$$x_1 + x_{n-k} + x_{n-k+1} + \cdots + x_n \le r^2/4 = k + \alpha$$

surely one can find $x'_{n-k}, x'_{n-k+1}, \ldots, x'_n$ such that

$$x_1 + x_{n-k} + x_{n-k+1} + \cdots + x_n = x'_{n-k} + x'_{n-k+1} + \cdots + x'_n$$

with $x'_j \ge x_j$ for $j = n-k, n-k+1, \ldots, n$. Now set $x_1^* = 0, x_j^* = x_j$ for $j = 2, 3, \ldots, n-k-1$ and $x_j^* = x'_j$ for $j = n-k, n-k+1, \ldots, n$. By the mean value theorem, there exist ζ_j's with $x_1^* = 0 \le \zeta_1 \le x_1$ and $x_j \le \zeta_j \le x_j^*$ for $j = 2, 3, \ldots, n$ such that

$$F(x_1^*, x_2^*, \ldots, x_n^*) - F(x_1, x_2, \ldots, x_n) = \sum_{j=1}^n G'(\zeta_j)(x_j^* - x_j).$$

Since $G(x)$ is a strictly convex function, we have

$$0 < G'(\zeta_1) < G'(\zeta_2) < \cdots < G'(\zeta_n).$$

$$\Delta(n, m) \le \sin \sqrt{\frac{\pi^2}{n} \left\lfloor \frac{(r_0^E)^2 (n, m)}{4} \right\rfloor + \frac{4}{n} \arcsin^2 \sqrt{\frac{(r_0^E)^2 (n, m)}{4} - \left\lfloor \frac{(r_0^E)^2 (n, m)}{4} \right\rfloor}}. \qquad (B2)$$
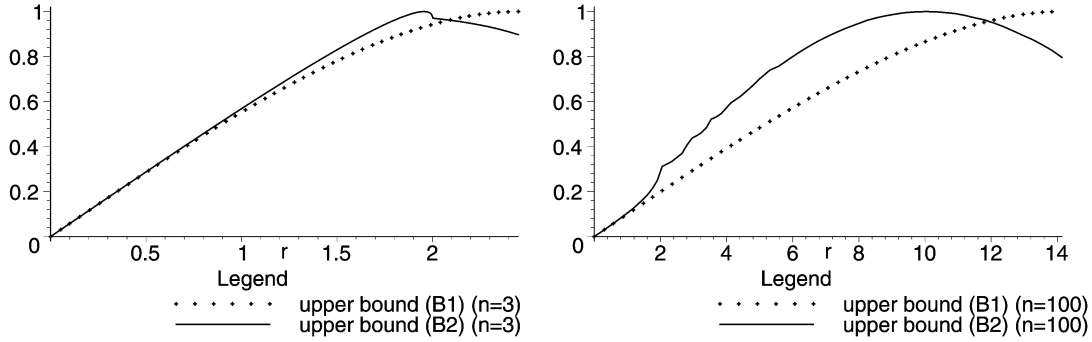
Fig. 1. The comparisons of two upper bounds as functions of $r = r_0^E(n, m)$ for $n = 3$ and $n = 100$.

Now

$$F(x_1^*, x_2^*, \ldots, x_n^*) - F(x_1, x_2, \ldots, x_n)$$
$$\geq G'(\zeta_2) \left(\sum_{j=2}^{n} (x_j^* - x_j) \right) - G'(\zeta_1)(x_1 - x_1^*)$$
$$= \left(G'(\zeta_2) - G'(\zeta_1) \right)(x_1 - x_1^*)$$
$$= (G'(\zeta_2) - G'(\zeta_1))x_1 > 0.$$

This contradicts the maximality of F at $(x_1, x_2, \ldots, x_n)$. Applying exactly the same analysis to $x_2, x_3, \ldots, x_{n-k-1}, x_{n-k}$, we deduce that $x_j = 0$ for $j = 2, 3, \ldots, n - k - 1$ and $x_{n-k} = \alpha$. So the upper bound of F can be given as

$$F(x_1, x_2, \ldots, x_n) \leq \left(k \frac{\pi^2}{4} + \arcsin^2(\sqrt{\alpha}) \right).$$

Take

$$g(r) = 2\sqrt{n} \arcsin(r/(2\sqrt{n}))$$

and

$$f(r) = 2\sqrt{k\pi^2/4 + \arcsin^2 \sqrt{\alpha}}$$

the corollary follows according to the previous theorem. □

Note that both upper bound (B1) and upper bound (B2) depend on $r_0^E(n, m)$. In Fig. 1, we plot both upper bounds as functions of $r_0^E(n, m)$ for 3 and 100 dimensions. One can see that if and only if $r_0^E(3, m) > 2.0881$, the upper bound (B2) is tighter than the upper bound (B1). While for the 100-dimensional case, the upper bound (B1) is tighter than the upper bound (B2) if and only if $r_0^E(100, m) > 11.9155$. In fact, it can be checked that asymptotically when n is large enough, upper bound (B2) is tighter than upper bound (B1) if and only if $r_0^E(n, m) > 1.1892\sqrt{n}$. For a packing problem on a manifold, alternatively, one can choose the neighborhood to be a small "ball" with Riemannian radius r. This will be our third approach to derive an upper bound for the diversity sum. For a particular $A \in U(n)$, let

$$U_r^R(n, A) = \{U \in U(n) | \text{dist}(U, A) \leq r\}.$$

Note that the constraint $\text{dist}(U, I) \leq r$ is equivalent to

$$\theta_1^2 + \theta_2^2 + \cdots + \theta_n^2 \leq r^2.$$

Therefore, we apply the same argument as in the proof of Theorem 2.3 and conclude that

$$V(U_r^R(n)) = \frac{\int\int_{D_1 \cap D_4} \prod_{j<k} |e^{i\theta_j} - e^{i\theta_k}|^2 d\theta_1 d\theta_2 \cdots d\theta_n}{\int\int_{D_1} \prod_{j<k} |e^{i\theta_j} - e^{i\theta_k}|^2 d\theta_1 d\theta_2 \cdots d\theta_n}$$
$$\times V(U(n))$$

where D_1 was defined in (2.5) and

$$D_4 := \left\{ (\theta_1, \theta_2, \ldots, \theta_n) | \sum_{j=1}^{n} \theta_j^2 \leq r^2 \right\}. \quad (2.16)$$

Instead of considering the Euclidean neighborhoods $U_r^E(n, A_1)$, $U_r^E(n, A_2), \ldots, U_r^E(n, A_m)$, we can consider the Riemannian neighborhood $U_r^R(n, A_1), U_r^R(n, A_2), \ldots, U_r^R(n, A_m)$. Utilizing the fact that the Euclidean distance $\|A_j - A_k\|$ and the Riemannian distance $\text{dist}(A_j, A_k)$ are related (compare with (2.15)

$$4n \arcsin^2(\|A_j - A_k\|/(2\sqrt{n})) \leq \text{dist}(A_j, A_k)$$

for any two unitary matrices A_j and A_k, we can derive the third upper bound (B3). The proof of the following theorem is very similar to the one of Theorem 2.8 and for the sake of brevity we omit it.

Theorem 2.10: Let D_1 and D_4 be defined as in (2.5) and (2.16) and assume $r_0^R(n, m)$ is the solution to the following equation (with variable r):

$$m \int\int_{D_1 \cap D_4} \prod_{j<k} |e^{i\theta_j} - e^{i\theta_k}|^2 d\theta_1 d\theta_2 \cdots d\theta_n$$
$$= \int\int_{D_1} \prod_{j<k} |e^{i\theta_j} - e^{i\theta_k}|^2 d\theta_1 d\theta_2 \cdots d\theta_n. \quad (2.17)$$

Then

$$\Delta(n, m) \leq \sin\left(\frac{r_0^R(n, m)}{\sqrt{n}} \right). \quad (B3)$$

Remark 2.11: Again suggested by Hassibi and similar to the analysis in Remark 2.5, we have

$$\int\int_{D_1 \cap D_4} \prod_{j<k} |e^{i\theta_j} - e^{i\theta_k}|^2 d\theta_1 d\theta_2 \cdots d\theta_n$$
$$= \frac{n!}{2\pi} \int \frac{1}{w}(e^{-iw} - 1) \det(F) dw$$

where F is defined in (2.10) with $f(\theta) = e^{i\frac{\theta^2}{r^2}w}$. Furthermore, by the following formula:

$$\int_{-\infty}^{\infty} e^{-ix^2} = (1/2 - 1/2i)\sqrt{2\pi}$$

one can explicitly calculate the entries of F as follows:

$$F_{jk} = (r/2 - r/2i)e^{ir^2(j-k)^2/4}\sqrt{2\pi}.$$

We gave three approaches to derive upper bounds for the diversity sum and hence also for the diversity product. All of them involve the
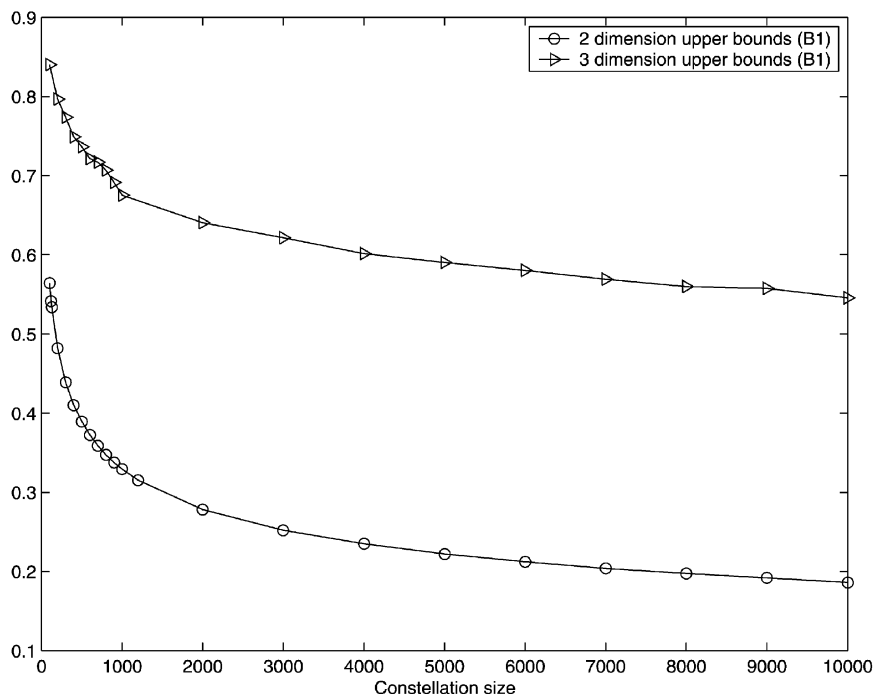
Fig. 2.   Upper bounds for two- and three-dimensional constellations.

TABLE I
FOR $n = 2$ THE TABLE COMPARES THE UPPER BOUNDS IN [14] WITH OUR NEW BOUNDS (B1) AND (B2)

m	24	48	64	80	100	120	128	1000
upper bounds in [14]	0.6746	0.6193	0.5969	0.5799	0.5632	0.5499	0.5452	
upper bound (B1)	0.7598	0.6603	0.6131	0.5932	0.5578	0.5425	0.5347	0.3270
upper bound (B2)	0.7794	0.6734	0.6235	0.6026	0.5654	0.5496	0.5415	0.3285

calculation of $r_0^E(n, m)$ or $r_0^R(n, m)$, which are the solutions of (2.8) and (2.17), respectively. Fortunately, we are dealing with finding a root of a monotone increasing function (recall that both $V(U_r^E(n, m))$ and $V(U_r^R(n, m))$ are monotone increasing functions with respect to r), the bisection method [2] will be highly effective to solve this kind of problem. Our numerical experiments for small-size constellations with small dimensions show that upper bound (B3) is looser than the first two upper bounds. However, when m goes to infinity, these three upper bounds give almost the same estimation. This makes sense because asymptotically the small balls look like a n^2-dimensional ball in Euclidean space. One can see the derived upper bounds for two- and three-dimensional constellations in Fig. 2.

We compare the derived upper bounds with the currently existing one presented in [14] (For computations, we refer to [7]). From Table I, one sees that the upper bounds ($n = 2$) of Liang and Xia [14] tend to be better when $m \leq 100$ and our bounds become tighter when $m \geq 100$.

One interesting fact about the limiting behavior of $\Delta(n, m)$ (when $m \to \infty$) is its connection to the Kepler problem [4]. Certainly one can use Kepler density [4] to obtain a tighter bound of the diversity sum asymptotically.

III. CONCLUSION AND FUTURE WORK

We presented three approaches to derive upper bounds for the diversity sum of unitary constellations of any dimension n and any size m. The derived bounds seem to improve the existing bounds when $n = 2$ and $m \geq 100$. When n is large, the exact computation of r_0^E is rather involved and hence it is also computationally difficult to compute the bounds (B1) and (B2). Nonetheless, it is our belief that the resulting upper bounds (B1) and (B2) become fairly tight as soon as m is sufficiently large.

It was pointed out that the resulting upper bounds also apply for the diversity product, although the bounds seem to be less tight in this situation. Future work may involve the derivation of a tighter upper bound analysis for the diversity product of unitary constellations using differential geometric means.

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Subchannel Allocation in Multiuser Multiple-Input–Multiple-Output Systems

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Abstract—Assuming perfect channel state information at the transmitter of a Gaussian broadcast channel, strategies are investigated on how to assign subchannels in frequency and space domain to each receiver aiming at a maximization of the sum rate transmitted over the channel. For the general sum capacity maximizing solution, which has recently been found, a method is proposed that transforms each of the resulting vector channels into a set of scalar channels. This makes possible to achieve capacity by simply using scalar coding and detection techniques. The high complexity involved in the computation of this optimum solution motivates the introduction of a novel suboptimum zero-forcing allocation strategy that directly results in a set of virtually decoupled scalar channels. Simulation results show that this technique tightly approaches the performance of the optimum solution, i.e., complexity reduction comes at almost no cost in terms of sum capacity. As the optimum solution, the zero-forcing allocation strategy applies to any number of transmit antennas, receive antennas and users.

Index Terms—Broadcast channel, multiuser multiple-input multiple-output (MIMO), orthogonal frequency division multiplexing (OFDM), successive encoding, sum capacity, zero-forcing.

I. INTRODUCTION

Increasing demand for broadband services calls for higher data rates in future wireless communication systems [1]. Data rates of several Mb/s for high mobility scenarios and up to 1 Gb/s in low mobility or static scenarios are expected in fourth generation systems. In the way to such transmission rates there are two major barriers to be overcome. The first is the scarcity of spectrum, which limits the amount of bandwidth available for transmission. The second is the wireless channel that severely distorts the signal due to multipath propagation.

The combination of multiple antennas and multicarrier technology seems to be ideal to achieve the expected rates under the mentioned constraints [2]. On the one hand, multiple-input multiple-output (MIMO) channels resulting from the use of multiple antennas at both transmitter and receiver show higher capacity than single-input–single-output (SISO) channels and this difference linearly grows for increasing transmit power. Thus, multiple antennas lead to higher spectral efficiency. On the other hand, multicarrier techniques, such as orthogonal frequency division multiplexing (OFDM), transform the frequency selective broadband channel into a set of nearly flat narrowband channels. As a result, distortion due to multipath is reduced and equalization at the receiver is greatly simplified.

In the work at hand, we consider the downlink of a wireless communication system with multiple antennas at the transmitter and the receivers and OFDM as transmission scheme. We assume that receivers